# FPGA IMPLEMENTATION OF AUTO TV TRACKING ALGORITHMS

# Manvendra Singh



TH EE/2002/M Si 64 f

Indian Institute of Technology Kanpur

# FPGA IMPLEMENTATION OF AUTO TV TRACKING ALGORITHMS

A Thesis Submitted in
Partial Fulfillment of the Requirements
for the Degree of

#### MASTER OF TECHNOLOGY

by

MANVENDRA SINGH

to the

DEPARTMENT OF

ELECTRICAL ENGINEERING

INDIAN INSTITUTE OF TECHNOLOGY

KANPUR

February 2002

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मुरुषोत्तम विकास मार्गातम भारतीय विकास विकास विकास सवान्ति क्राउं ते विकास विकास





This is to certify that the work containing in the thesis entitled "FPGA IMPLEMENTATION OF AUTO TV TRACKING ALGORITHMS" by MANVENDRA SINGH has been carried out under my supervision and that this work has not been submitted elsewhere for the award of a degree.

Dr. BAQUER MAZHARI

Assistant Professor,

Department of Electrical Engineering,

Indian Institute of Technology, Kanpur.

February 2002.

**ACKNOWLEDGMENT** 

I express my sincere and deep gratitude to my thesis supervisor Dr." B.

Mazhari for his inspiring guidance and encouragement at every stage of this work.

He provided me with numerous valuable, innovative and constructive suggestions,

without which this work would not have been possible. No words can aptly

describe his patience and understanding. I would also like to thank all those

instructors who have taught me various courses during the academic program

with special mention of Dr. R. Sharan who has been throughout a source of

encouragement and guidance.

I am thankful to Instrument Research & Development Establishment

(I.R.D.E.) of D.R.D.O. for providing me this wonderful opportunity. The support from

my colleagues of this establishment has been very valuable. The students who

have done this course with me have also been extremely helpful. I am grateful to

all of them.

I am extremely grateful to my parents, relatives and friends for their love,

support, blessings and encouragement. My wife Manju deserves a special

mention here for her patience and support.

February 2002

Manvendra Singh

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### **ABSTRACT**

Real-time target tracking in video sequences is an essential component of any Electro-optical system. The auto TV tracker computes the target position in successive frames with respect to the center of the field-of-view (FOV) using image tracking techniques. In the present work centriod and correlation based tracking algorithms have been discussed, designed and implemented into FPGAs using hardware description languages (i.e. VHDL). The centroid algorithm has been designed using loop pipelining method for real-time application. This design approach has been successfully simulated and implemented into Xilinx's FPGA XC4008E. For the implementation of correlation based tracking algorithm, three-memory approach has been used to cater the real-time requirement. The correlation based tracking uses the Sequential Similarity Detection Algorithm (SSDA) for the calculation of match point as it require less hardware than the classical algorithm. This technique has been designed, simulated and successfully implemented into Xilinx's FPGA XC4020E. The advantage of using FPGA implementation is that it can work as a stand-alone system and may be integrated with other system by using the standard communication protocol. This design approach has given the system, much-needed reconfiguribilty to suit the customer's requirement. It has also resulted in reduction of PCB fabrication time since the hardware remains essentially the same although the system design may change.

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# CHAPTER 1

# An Introduction to Auto TV Tracker

#### 1.1 Introduction:

A typical TV tracker system consists a TV camera, monitor and servo mechanism. The operator analyzes the image coming from camera. If the operator sees some target of interest on monitor then he/she will follow the target through the manual control switch (i.e. joystick).

This type of TV systems suffered from a significant number of problems such as:

- The necessity of having an operator to monitor the video signal, compromises overall system capabilities.
- The TV system is "dead weight" when an operator is not monitoring the video.

To ease the operator work and incorporate some intelligence in the system, an automatic TV tracker is used. The Auto TV tracker process the contrast information contained in the video signal obtained from TV camera (either CCD or Thermal Imager) and generates the horizontal and vertical error signals using the tracking algorithms. These error signals are further used to drive servo system on which the sensors are mounted.

#### 1.1.1 Need of TV Tracker

The auto TV tracker has wide range of applications in defense systems. In defense system it is used in almost all the systems, where the video monitor is provided to the operator for the analysis of scene. It is used in land defense system such as Main Battle Tank (MBT) for the alignment of Gunner's main-sight to give the accurate hitting of target. The TV tracker is also used in

the missile-seeker heads, Remotely Piloted Vehicles (RPV's) as well as in conventional aircraft and helicopters. It is widely used in the test ranges to track the missile path trajectory during the missile test fire. If suitable range data is provided through Laser Range Finder (LRF) then auto TV tracker alongwith the LRF could be used in automatic Fire Control System (FCS).

The auto TV tracker is a key component in the Electro-Optical Fire Control Fire Control System (EOFCS) for naval ships. Where it provides the continuous azimuth (Az) and elevation (El) position error data of the target of interest based upon the video image supplied by the EOFCS sensors. Using these positional data provided by the TV tracker, FCS compute the fire trajectory that finally goes to the Gun Control System for accurate hitting.

#### 1.1.2 Working principle:

The Typical auto TV tracker system configuration is shown in the fig.(1.1). An auto TV tracker receives a video input, locks onto a target selected by the operator within the TV field-of-view (FOV) and outputs error signals, describing the displacement of the target from the center of the field-of-view. A TV formatted system [either CCD, Thermal Imager (FLIR) etc.] produces a video signal in a sequential manner. The exact scanning format depends upon the sensor design (for this case it is of CCIR-B slandered i.e. 625 lines with 20 ms frame rate). In most sensors scanning begins at the upper left, sweeps to the right, continuing line-by-line down through the scene till the end of frame. By detecting the beginning of each line (i.e. horizontal sync), the end of the frame (i.e. vertical sync), and by using a high speed clock, the video signal may be broken into an X-Y coordinate matrix (i.e. Az, El) of TV lines. Now by choosing the suitable tracking algorithm such as (Edge, Centroid or Correlation, the auto TV tracker computes the azimuth and elevation position signal derived from the video contrast information of the selected target within the sensor field-of-view.

FIG 1: SYSTEM CONFIGURATION OF AUTO TV TRACKER

These positional azimuth and elevation error signals are supplied to servo system, which brings the target again in the center field-of-view. The target position has to be calculated and updated in each frame (i.e. every 20 ms) to give the real-time effect. Hence the hardware used to develop this system should be faster & algorithms should be optimized enough to cater the real time challenge.

#### 1.1.1 TV Tracker Specifications:

The typical auto TV tracker specifications are as follows:

•	Video input	Composite video: CCIR-B std.
		(625 lines, 50 HZ frame rate)

	•	Track gate mode	Manual & Auto
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•	Track gate	Center of screen (default) &
	position	Movable to any FOV position

•	Tracking algorithm	Centroid & Correlation
		(Selectable to the operator)

		(for correlation mode tracking)
		Search area 64 pixels X 64 lines/ frame
		target gate size 16 pixels X 16 lines
		(for centroid mode tracking)
	Size	Maximum 64 pixels X 64 lines/ frame
•	Track gate	Minimum 24 pixels X 24 lines

•	Tracking error	less than	1/2 pixel
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•	Minimum Target	less than 3 TV/ line = 80%
	Contrast	greater than 4 TV/ line = 25%

Minimum Target Horizontal = 0.8 % FOV
 Size Vertical = 0.4 % FOV

• Error output 8 bit (azimuth & elevation)

#### 1.1.2 Tracking Algorithms:

There are three algorithms widely used in the contrast based Auto TV tracker given as:

- Edge tracking
- Centroid tracking
- Correlation tracking

The Edge tracker seeks and acquires the first selected edge (i.e. right, left, top or bottom) of the target of interest, occurring within the gate, and maintains track on the edge. Smaller electronic gates may be used in edge tracking to gate out a major portion of the target and allow processing of only the selected track space point area. Thus, the gate of the edge tracker may be used to reject unwanted background signals. This is a significant advantage when tracking small target against a cluttered background. When the selected edge of the target occurs, horizontal and vertical position counters are sampled and the position get stored. These signals represent the position of the target edge, with respect to the internally generated reticle, or offset track point. These data are supplied as an output in digital form to the servo mechanism, which in turn follows the selected edge.

Two difficulties may arise in edge tracking procedure

- (i). Since the differentiation is typically restricted to gated horizontal signals (i.e. differentiation of gated video in each TV lines), target that are of large extent than the track window and parallel to horizontal datum provides no track information.
  - (ii). Since the tracker must contain a corner within the track window

it continues to indicate an error until a corner is reached. An example of this effect is noted when the tracker locks on to a picture of a road or cloud edge and is committed to follow this edge until the track window reaches its excursion limits.

The centroid tracking is commonly used for air and naval surface targets, because these tend to be bounded targets. The centroid algorithm is so called because it finds the centroid of the target, and then designates the target aim point. The video signal from the TV camera contains a complete description of the entire field of view (FOV) of the image (target size, shape, shades of gray etc.) within the limit of the capabilities of the camera. The entire FOV of the camera appears on the TV monitor but an adjustable gate reduces the sample area of camera FOV to be used by the TV tracker. In general, the centroid mode of tracking is preferred due to its inherent immunity to noise and false alarm.

Correlation TV tracking is a process of acquiring target position information using a reference image and comparing it (each TV field) to the real time video in a defined "search window" within the FOV. The Correlation mode of tracking is used to track a target of defined shape and will maintain a track on a single feature of an extended target against a cluttered background. It maintains lock on a target, which is temporarily obscured.

The brief criteria of the selection of above discussed algorithms depends on scenario and are shown in the table 1.

CONDITIONS	TRACKING ALGORITHMS
Dynamic Target	Centroid
Bounded Target	Centroid
Varying background	Correlation
Target Contrast (High)	Centroid
Target Contrast (Low)	Correlation
Target Size (Large)	Centroid
Varying Target Shape	Centroid

#### 1.1.3 Performance Characteristics of TV Tracker:

To accurately define the performance of TV camera, TV tracker and its servo system interface, the parameters such as range-vs-target size, target contrast and tracking bandwidth must be considered.

Range-vs-target size is most easily defined in terms of target television lines (TVL) per field as given in equation (i)

This equation gives the following conclusions:

- If the above equation results in TVL/ field = 1, the target is theoretically detectable by the TV systems.
- If the equation results in 2≤TVL/ field ≤ 3, the target is trackable but represents a generally poor response from the sensor.
- If the equation results in TVL/ field ≥ 4, the target is easily tracked and the sensor response is near 100 %.

In addition to target size, the contrast of the target is an important consideration. The minimum target contrast may best be explained in terms of signal-to-noise (S/N) ratio for electro-optical sensor system. The equation (ii) gives the S/N (peak-signal-to-RMS noise) for such system.

$$S/N = \frac{(C) (\alpha) [MTF (L)] [MTF (S)] (G)}{Noise} - (ii)$$

where

C = true contrast of the target to background in the spectral region of the sensor.

 $\alpha$  = atmospheric attenuation in the spectral region of the sensor.

MTF(L) = modulation transfer function of the lens.

MTF(S) = modulation transfer function of the sensor.

G = sensor response to the signal present at the image plane.

This equation (ii) gives the following conclusions:

- If the equation results in a S/ N ≥ 2, the target may be tracked by a centroid tracker.
- If the equation results in a S/  $N \ge 4$ , the target may be tracked by an edge tracker.

#### 1.1.6 Pre-processing:

Since the quality of video signal coming from the CCD or Thermal Imager heavily depends on sensor quality and the scene environment, so before the Auto TV tracker unit performs the selected tracking algorithms (i.e. Centroid or Correlation), pre-processing is required. The pre-processor makes the video signal directly useful to TV tracker. Further, pre-processor enhances track quality and incorporate some intelligence in tracker unit. The following processing is required to make the video signal useful to TV tracker:

- Pre-processing for noise suppression.
- Image processing for separation of the target from heavy clutter background.
- Image enhancement and histogram equalization (for contrast enhancement).

On completion of the above described processing the video signal will be free from noise as well as it will have acceptable contrast for TV tracker.

To provide better synchronization and ease to the operator an Auto TV tracker requires following additional circuit along with the tracking algorithms:

- Sync separator
- Threshold mode and threshold circuit
- Track Gate mode selection and Gate size

#### Sync Separator:

The sync separator circuit used to separate horizontal synchronization signal (i.e. H Sync, which gives the start and end of one line) and vertical synchronization signal (i.e. V Sync, which gives the start and end of frame) from the composite video signal coming through image sensor. These separated H sync & V sync information is applied to Auto TV tracker circuits for the synchronization of timing and control functions. Fig. (1.2) shows the circuit diagram for separation of sync-signal from the composite video.

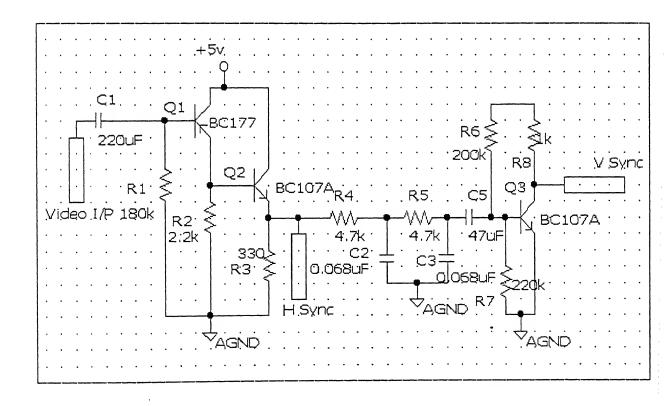


Fig. (1.2): Schematic Diagram of Snyc Separator Circuit

#### Threshold control:

The video signal coming form sensors are analog in nature. For the implementation of tracking algorithms (which are in digital format), video signal has to be converted in digital format. The process of converting analog video signal to binary information is called thresholding.

In this control, the operator selects a video intensity threshold, which is just below the dominant intensity of the target. The operator has to select white hot/black hot target from analysis of the scene. It means, whether the target has the white contrast against the black background or black contrast information against the white background. Fig. (1.3) shows scheme for white hot/black hot threshold waveform.

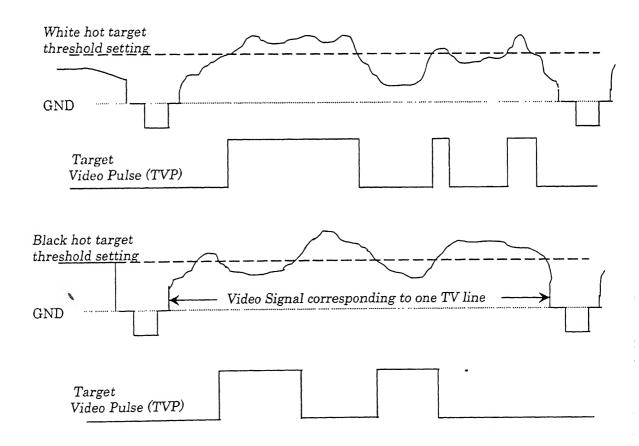


Fig. (1.3): Waveform of Thresholding white and black hot target data

results. Chapter 3 gives the implementation and algorithmic details of the correlation tracking. Chapter 3 also describes the modules required for the implementation of correlation tracking alongwith the simulation and synthesis report. Appendix-I gives the implementation reports such as pad and map reports of the centroid and correlation algorithm implemented in respective FPGAs. Appendix-II describes the Xilinx's FPGA configuration scheme and pin details of the targeted FPGA devices used for the implementation of above proposed algorithms.

#### Track Gate:

Track gate is used to eliminate information other than the target of interest selected by the operator. For this purpose a track gate is positioned at the center of the screen. A switch is provided to the operator for the selection of Auto or Manual track gate size mode. This option operates in centroid mode of tracking only.

In the manual mode the operator changes the track gate size (through the UP/ DOWN switch provided in operator control console), depending upon the size of the target. This is done to ensure single target tracking.

In auto mode the track gate automatically increases or decreases in size with target variations, to maintain a constant GATE/ TARGET area relationship.

## 1.2 Objective of Thesis:

The objective of the thesis is to design and implement centroid & correlation based TV tracking algorithms in on FPGA for real-time application. Since the auto TV tracker is used in real-time applications, the timing constraint has to be fulfilled (i.e. the computation should be completed within one frame time). The target FPGAs for the above algorithms are XC 4008E and XC 4020E with the capacity of 8000 & 20000 gates. The maximum speed of operation of the target FPGA is 80 MHz. For the FPGA design implementation Xilinx's XACT tool has been used.

## 1.3 Organization of the work:

The requirements of auto TV tracker and it's algorithmic details were discussed in the present chapter. Chapter 2 describes the actual FPGA design methodology for the implementation. This chapter also describes the details of the centroid tracking algorithm to be implemented alongwith the

# **CHAPTER 2**

# **Centroid Based Tracking Algorithm**

#### 2.1 Background:

The centroid based tracker is an image integrating tracker used mainly for tracking the bounded target and advantageous in the scenario where the aspect ratio of target is continuously changing e.g. Arial targets. The video signal from the TV camera (CCD or Thermal Imager) contains the complete description of the entire field-of-view (FOV). The entire FOV of the camera appears on the monitor but an adjustable track gate reduces the sample area of the FOV to be used by the tracker to compute the error signal in azimuth and elevation direction. Centroid based tracker integrates the video within the track gate by partitioning the active area into a matrix as shown in fig (2.1).

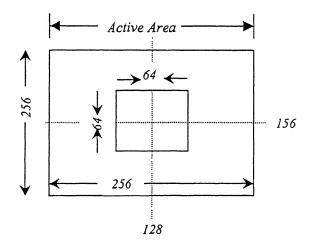


Fig.(2.1): Scheme of Partitioning the active area

The Horizontal sync in the CCIR-B standard video format has the period of 15.625 KHz. So if 4 MHz clock is chosen for scanning then the active area can be

defined in terms of (4 MHz/ 15.625 KHz) = 256 pixels. The track gate size is shown in fig. (2.1) is for 64 pixels/64 lines.

For computation of centroid of the target within the track gate, the video signal is digitized using threshold scheme to get the binary image as shown in fig. (2.2).

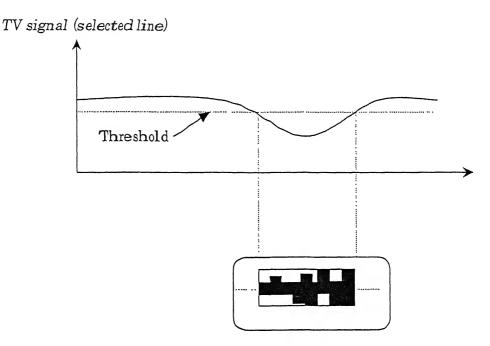


Fig. (2.2): TV picture (binary target representation)

The above figure shows the binary target representation of the video signal corresponding to one line. The video information is digitized in successive lines to generate target video pulses (TVP). The tracker integrates the horizontal and vertical coordinate data of the digitized video signal to compute the centroid of the target within the track gate using the equation (i) and (ii) given below:

$$X_{CENTROID} = \frac{\sum_{i=0}^{m} X_i}{\sum_{i=0}^{n} N_i} \qquad ---- \text{(i)}$$

$$Y_{CENTROID} = \frac{\sum_{i=0}^{m} Y_i}{\sum_{i=0}^{n} N_j} \qquad ---- \text{(ii)}$$

Where

N = data point

 $X_i = X$ - position of the target

 $Y_i = Y_i$  position of the target

n = total number of data points

m = total number of position coordinates

These centroid signal are use to compute the error signal with respect to the center of screen and being send to servo system, which will bring the target again in the center of the monitor in next frame.

#### 2.2 FPGA Based System Design:

A typical flow diagram for the FPGA based system design methodology has been shown in the fig. (2.3). The design starts with system specification, which is converted into the modules and sub-modules. These modules are initially designed, synthesized and simulated functionally using Hardware Description Language (e.g. VHDL). These modules after the simulations are interconnected according to the specification. After the interconnection, the complete design is synthesized and simulated for the particular FPGA device. The gate level synthesis gives the actual gate density required for the design. After the layout synthesis a bit file is generated for the target FPGA. This bit file is used for the configuration of on board FPGA device.

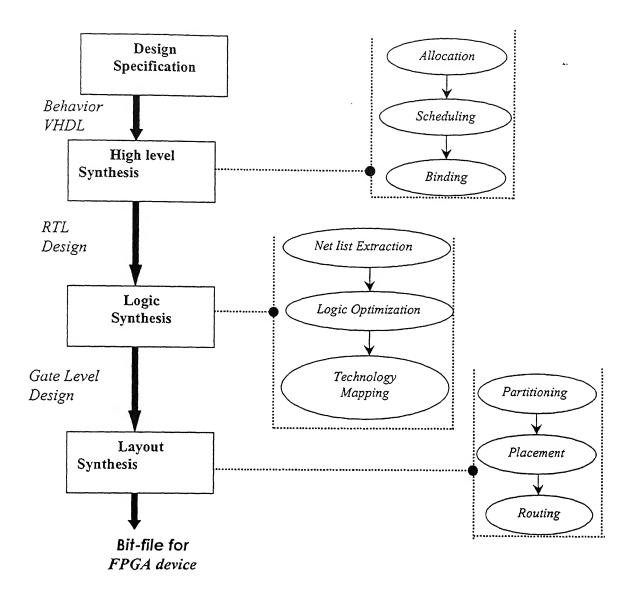


Fig. (2.3): Flow diagram for FPGA based design methodology

# 2.3 Design of Centroid Based Tracking Algorithm:

For actual implementation of the centroid algorithm, scheme is shown in fig. (2.4). The video signal is first digitized using the video threshold circuit with the help of reference voltage so that the target contrast level should be above the threshold level. This will generate the binary image of the target of interest. To reject the information other than the target, the track gate is generated with the two selectable modes. One is Manual mode and second is auto gate mode. The operator can select either of these modes through control switch provided on

operator control console. The track gate alongwith binary information is processed to generate target video pulse (TVP). Basically TVP generator is AND gate which give the binary information within the track gate. The TVP within track gate goes to the computational logic block, which computes centroid of the target using eq. (i) & (ii) with horizontal and vertical sync as a control signal.

All logic except the threshold is digital in nature. In the present work, implementation of computational logic is modeled in the VHDL and finally fused in the Xilinx FPGA XC 4008E.

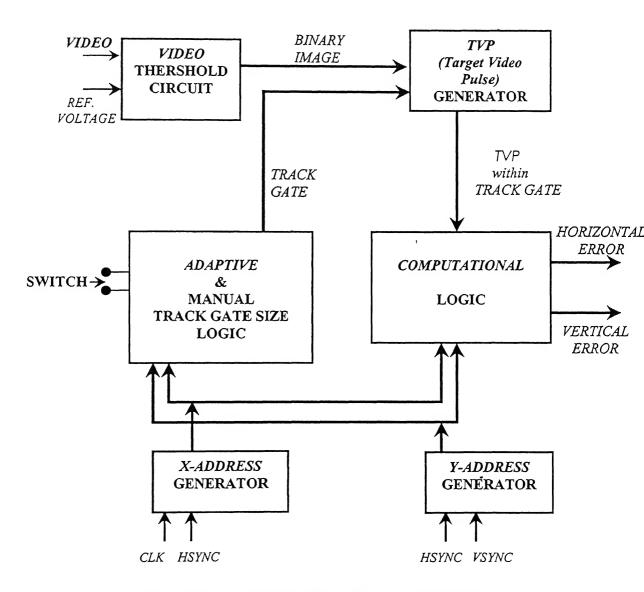


Fig.(2.4): Scheme of Centroid Tracking Implementation

#### 2.3.1 Generation of Manual Track Gate:

The scheme for the generation of manual track gate size is shown in fig. (2.5).

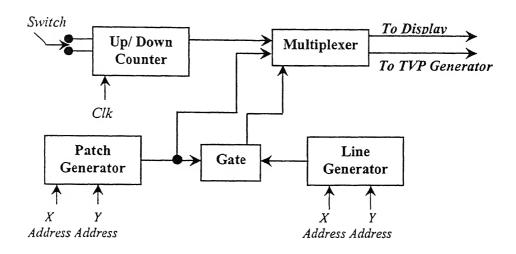


Fig. (2.5): Scheme for Manual Track Gate Size

In the manual mode track gate size the operator has the flexibility to increase or decrease the track gate size depending upon the scenario, within eight steps from maximum (64 pixels X 64 lines) to minimum (24 pixels X 24 lines). For this purpose a switch is provided. The patch of above sizes and corresponding lines are generated using horizontal and vertical address generator in separate modules in VHDL. These lines and patches are passed through the AND gate. These gated lines are used for display purpose, which gives the visual effect to operator (i.e. the target of interest is confined to the area covered by line pattern). The output of the AND gate alongwith the patch generated goes to the multiplexer. The multiplexer output ensures that the same patch pattern (to be used as a control signals by the computational block & TVP generator) alongwith line pattern goes to the display.

#### 2.3.2 Generation of Automatic Track Gate:

The scheme for the generation of automatic track gate size is shown in fig. (2.6).

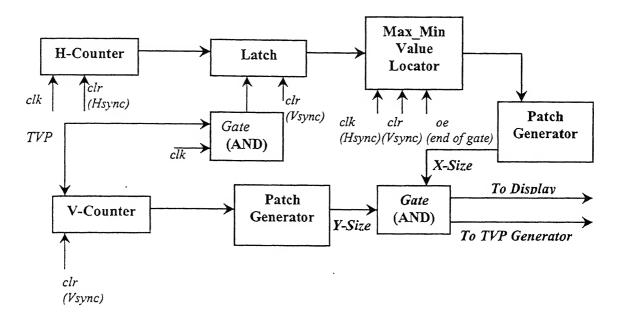


Fig. (2.6): Scheme for Auto Track Gate Size

In auto track gate size mode, the track gate automatically increases or decreases in size with target variations, to maintain a constant GATE/ TARGET area relationship. This will lead to single target tracking. In this mode the initial track gate size will have its default value and in the next frame onwards it will confined the as per the target size. The H-counter will generate the address using 4 MHz clock and h-sync. The TVP inside the clock is used as a clock to the latch, which will store the corresponding address of the target. The Max\_Min value locator is used to store the maximum and minimum value of the address latched. Those will be the lower and upper limits of the target size. The patch generator will generate the patch in the X-direction using these limits. For the Y-direction the vertical counter will count the no. of TVP in a frame and will generate the patch. Finally these patches are used to generate the control signal for computational logic and lines for the display.

#### 2.3.3 Computational Logic:

The calculation of the centroid of the target is computationally intensive. Since the tracking is a real-time operation hence the hardware should be fast enough and logic should be optimized to cater the requirement. The implementation scheme of computational logic is shown in fig. (2.7).

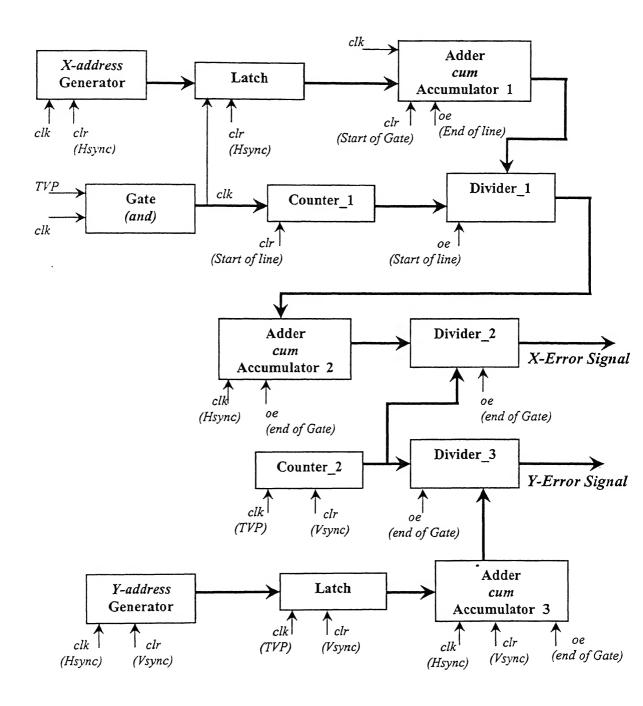


Fig. (2.7): Scheme for Implementation of Computational Logic

The X- address generator will give the continuous address as per the partitioning scheme suggested in fig. (2.1). The latch will store the corresponding address of the target using the clocked target video pulse (TVP) as a clock. These latched addresses will get added and accumulated for a single line. At the end of the track gate line, the accumulated value will be applied to the divider\_1 logic block, which is the dividend. At the same time the counter\_1 block will count the number of clock pulses inside the TVP and fed to the divider\_1, which is the divisor. The divider\_1 block gives the quotient as an integer. The quotient will be fed to the accumulator\_2 with the rising edge of the h-sync. This operation will continue for each line upto the last line of the track gate. Finally the accumulator\_2 value will be divided by the number of TVP present in one frame. This will be the XCENTROID or the horizontal error signal of the target.

For the calculation of the YCENTROID or the vertical error signal of the target, the Y-address generator will give the vertical address using h-sync as a clock and v-sync as a clear signal. The latch will store the corresponding address with the every rising edge of the TVP. The accumulator\_3 will add & accumulate all the latched addresses and fed to the divider\_3 block, which will divide the accumulated value by the number of TVP counted by the counter\_3. This is the Y-error signal of the target of interest.

Fig. (2.8) shows the waveform of the control signal such as target video pulse (TVP), clocked TVP (i.e. the clock inside the TVP) etc. This figure also shows the complexity involved in the calculation of the centroid of the target. For the case of 64 X 64 track gate size, time left for the calculation is 24  $\mu$ sec. for each line. So the conventional discrete IC's can't do the calculation that fast, but the FPGA (with 1-3ns delay) can solve the purpose. The key component in this design is the divider logic, which takes most of the time. The design of the divider logic is discussed in the next section. To avoid any timing conflict the control signal is modified to carry the accumulated value till the start of next track gate line. That is *loop pipelining* is used, which will increase the division time from 24  $\mu$ sec to (24  $\mu$ sec) = 48  $\mu$ sec.

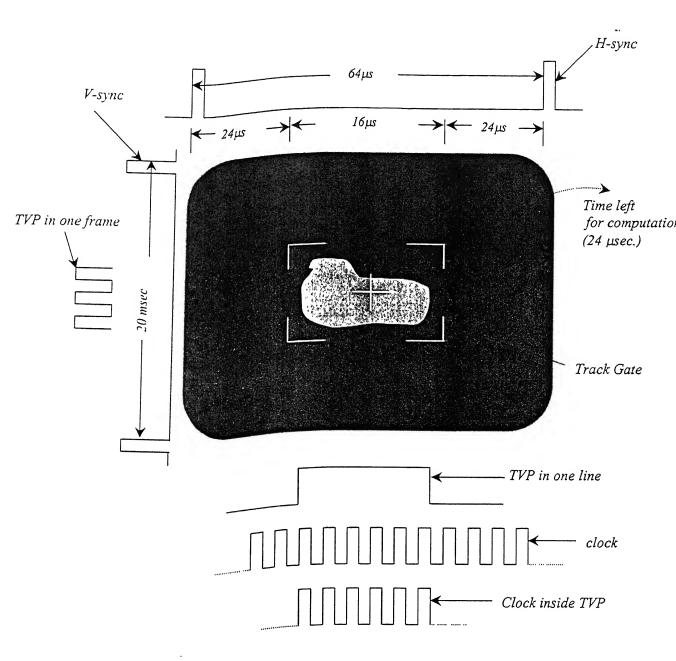


Fig. (2.8): Waveform of Control signals used in the Computational Logic

### Design of divider:

The logic used for the design of divider is shown in fig. (2.9). The components used for this logic are comparator & substractor, multiplexer and counter. For the first clock cycle MUX will send the input A (i.e. dividend) and it will get compare with the input B (i.e. divisor).

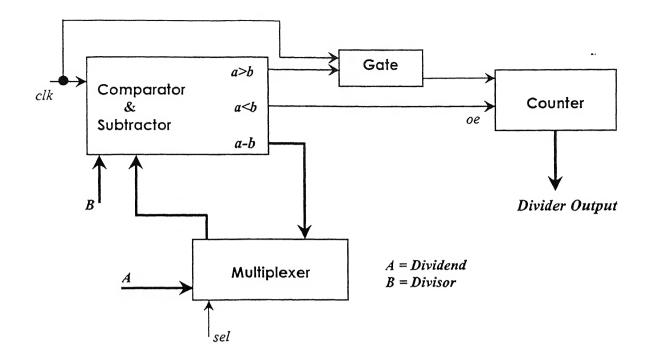


Fig. (2.9): Scheme for Divider Logic

If this value is less than A, then this block will send the subtracted value to MUX as its second input. Simultaneously the a>b will be high. In second clock onwards the MUX will send its second input and similar operation will be repeated till the a<b />
b is high. As long as the output a>b is high, the clock signal gated with a>b serves as a clock to counter. The counter counts the no. of clock pulses and this will be the output of divider (i.e. quotient). The divider shows its output only when the a<b />
b is high (i.e. output enable).

#### 2.4 Results and Discussion:

The divider logic have been simulated and synthesized first to check the requirement. The figure (2.10) shows the simulation result of the divider logic. Since the designed divider is clock cycle dependent, hence it can be best analyzed by the worst-case inputs. For the case of 64 X 64 track gate size, the starting address of the line is 96 and ends at 160 (according to the partitioning

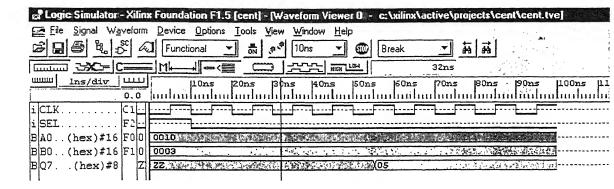


Fig.(2.10): Simulation result of the divider implemented

scheme suggested above). So if all the pixels are active i.e. the whole line of the track gate is covered by the target then the divider requires the 8320/64 = 130 clock cycle to complete the division operation. With the clock of 4 MHz it will take 32.5  $\mu$ sec, but using loop pipelining the available time is 48  $\mu$ sec. For the worst case, when only the last pixel of the track gate line i.e.  $160^{th}$  pixel is active, then the divider will require the 160/1 = 160 clock cycle to complete the division operation. It means it will take 40  $\mu$ sec. So the 8  $\mu$ sec time is left to cater the interconnect delay. Fig. (2.7) shows the divider output (05) $\mu$  with the inputs (0010) $\mu$  and (0003) $\mu$ .

The fig. (2.11) shows the simulation results of the computational logic implemented to calculate the centroid of the target. The target video pulse (TVP) is been detected at the addresses 03, 04 and 05, which are added and accumulated. The counter counts the no. of clock cycle inside the TVP i.e. 03. The divider gives the final X-error as 04. Similarly for the Y-error the vertical address 04, 05 and 06 are added and divided by the no. of the TVP in a frame i.e. 03. So the Y-error is 05 as shown in the simulation result.

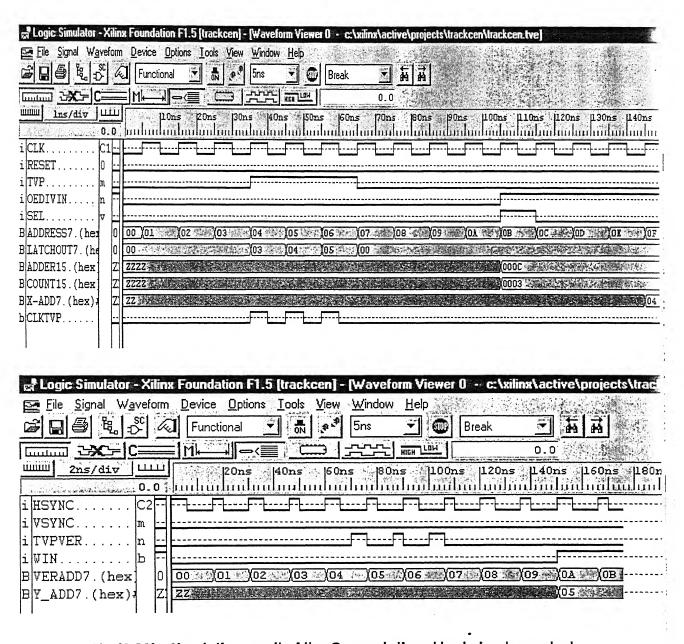


Fig.(2.11): Simulation result of the Computational logic implemented

For the worst case, divider requires 160 clock cycles. Since the available time for the division is 48  $\mu$ sec. The required operating frequency will be 48/160 = 0.3 $\mu$ sec = 3.3 MHz. It means that the system can use the clock frequency of 4MHz. The point to be noted that if the loop pipelining is not being used then the minimum required frequency would have been 6.6 MHz, as the available time for the division reduces to 24  $\mu$ sec.

The comprehensive synthesis and implementation result of the centroid tracking algorithm into FPGA XC4008E has given below:

Number of CLBs : 233 out of 324

CLB Flip Flops : 190 4 input LUTs : 115

3 input LUTs : 54 (15 used as route-through)

Number of bonded IOBs : 25 out of 61

IOB Flops : 0
IOB Latches : 0
Total equivalent gate count for design : 6425
Additional JTAG gate count for IOBs : 1200

The detailed map & pad reports are attached in the appendix-1.

# **CHAPTER 3**

# **Correlation Based Tracking Algorithm**

#### 3.1 Background:

Finding the location of a reference image within the field of view (FOV) of a video sensor in real time is considered using a correlation technique. Real time means that the reference image location within the live video image FOV is found for each field of the video sensor. This requirement prevents not only a buildup of data requiring large memory but also does place severe speed requirements on the algorithm chosen. For a standard image sensor (CCD of thermal imager) with a frame rate of 25 Hz, 2:1 interlace, and corresponding field rate of 50Hz, each match point must be computed in 1/50~Hz = 20ms. The correlation tracking is an effective method of targets that do not have the salient characteristics. The correlator essentially computes a surface or map of coefficient, which are a measure of similarity of the reference to all portion of the scene. The location of the highest match is used to update the aim-point of the sensor.

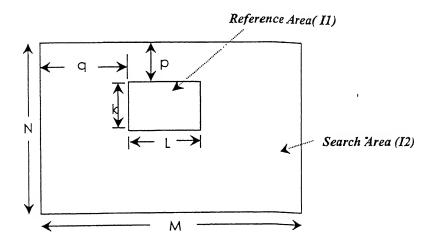


Fig. (3.1): Geometry of correlating reference image with search image

Figure (3.1) shows the geometry of correlating  $K \times L$  reference image (11) with  $N \times M$  sensor live image (12). Both the images are digitized and quantized into pixels. Additional preprocessing of one or both images will be necessary if the assumptions described in the chapter1 are not met.

Two algorithms are known for the implementation of correlation tracking technique: one which is the discrete version of the classical definition of correlation, and other which is related to this classical definition.

The classical approach to the problem of determining where two signals match is correlation. Consider two functions  $f_1(t)$  and  $f_2(t)$ , the correlation integral can be defined as equation (i)

$$C(\gamma) = \int_{-\infty}^{\infty} f_1(t) f_2(t+\gamma) dt \quad \text{-----} \quad \text{eq. (i)}$$

where  $\gamma$  is allowed to take the values  $-\infty$  to  $+\infty$ . The value of  $\gamma$  maximizes  $C(\gamma)$  in equation (i) is the correlation peak and is defined to be the match point between the two signals. It is obvious that determining the correlation peak consists of multiplying one signal by the other signal, shifted by  $\gamma$ , and then evaluating the area under resulting curve.

If the both signals are sampled and held, then equation (i) can be approximated by the following equation:

$$C(p) = T \sum_{n=-\infty}^{\infty} f_1(n) f_2(n+p)$$
 -----eq. (ii)

Where  $f_1(n) = f_1(nT)$  and  $f_2(n+p) = f_2(nT+pT)$  and T is the sampling interval. As T becomes very small, the result in eq. (ii) approaches that of eq. (i) where the p which maximizes C(p) is defined as the correlation point.

The digitized video image (I2) is represented by an N x M array of pixels as shown in fig.(3.1). The values of N and M are fixed by the choice of sample rate and number of TV lines per frame. For our design it is  $64 \times 64$ . The reference image (I1) is represented by a K x L array of pixels. For our design it is  $16 \times 16$ . The p and q dimensions in fig. (3.1) give the vertical and horizontal position,

respectively, of the reference in the video image. These indices start in the upper left corner of 12, where p = q = 0. One possible two-dimensional discrete correlation algorithms (often referred as the direct method) is given by

where R(p, q) is the correlation function, and the division by KL is the scaling tactor. In eq. (III) to each possible value of pland q, F(p, q) is obtained by summary in a product of line to expressionaling pixels in the resonance R(p, q) is the image registration or match point.

Another algorithms for determining the similarity or the amount of registration between two images is known as a sequential similarity detection algorithm (SSDA). The SSDA registration surface is computed by adding the KL absolute values of the difference between corresponding I1 and I2 pixels for each value of p and q as given by equation (iv).

$$E(p,q) = \sum_{n=1}^{K} \sum_{m=1}^{l} |I1(n,m) - I2(n+p,m+q)| ------ eq. \text{ (iv)}$$
 for  $0 \le p \le N$ -K,  $0 \le q \le M$ -L

The registration point (denoted by p\* and q\*), is the value of p and q for which E is minimum.

For the implementation of correlation tracking technique in this work, SSDA method has been chosen because of the following reasons:

 The SSDA method implementation requires less hardware since it uses subtraction rather than multiplication (used in direct method). Hence the implementation of direct method would require the tracker generally to have extensive hardware/ software capability and normally cannot be

- performed in real-time, which is a basic requirement for electro-optical tracking systems.
- Since the SSDA method uses the absolute difference of the search and reference data, hence it can perfectly register the white target against the black background as well as the black target against the white background. But the direct method cannot have this flexibility. The operator doesn't have to change the polarity of the target in correlation mode of tracking, when SSDA method is used.

### 3.2 System Configuration of Correlation Tracking Algorithm:

The scheme of the system configuration is shown in fig. (3.2). The tracker accepts the video signal from image sensor either CCD or Thermal Imager.

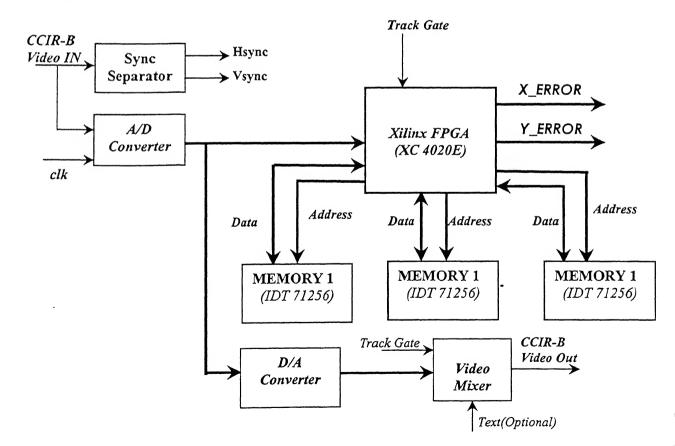


Fig. (3.2): System Configuration of Correlation based Tracking

The sync separator gives the h-sync and v-sync from the video signal received, as per circuit shown in fig . (1.2). The A/D converter (using AD9048) digitizes the analog video information of the scene into 8-bit data format. These ADC data are used in the computation of match point. In this system three memory approach has been used to cater the real-time requirement. The live image (64 x 64) will be stored first in the memory1 and the next frame will be stored in memory 2( 64 x 64). At the start of the third frame the data form the center of the first memory (16 X 16 size) will be read and it is the reference data. The memory 2 will be treated as the search data. Simultaneous the memory 3 will store the third frame. Computation logic (SSDA) will compute the best match point and send the error signals to the servo system, which will bring the reference image portion again to the center of FOV. In the fourth frame the memory 2 will give the reference data, memory 3 will give the search data (track data), memory 1 will be in the writing mode and vice-versa. This scheme has a latency of one frame, but computes the position of the target in a one frame (i.e. in real-time).

The computational logic, address scheme for the reference & search memory and control logic have been designed in VHDL and implemented in FPGA (XC4020E).

## 3.3 FPGA implementation of Correlation Based Tracking:

Fig. (3.3) shows the scheme of implementation of correlation tracking algorithm using FPGA. The design has been done by using the various modules such as main controller, address generator, interconnect logic and computational logic.

Main Controller: The main controller ensures that the operations with respect to the three memories should change accordingly with every three-frame pulse.

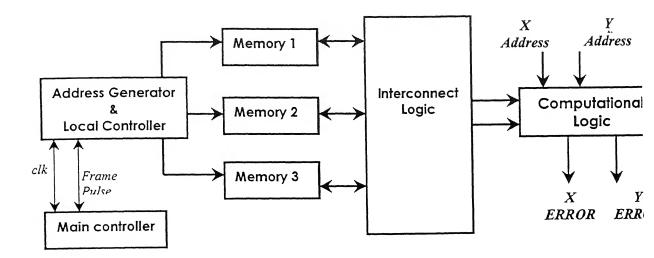


Fig. (3.3): FPGA Implementation of Correlation Tracking Algorithm

Address Generator: The address generator gives the writing and reading address to the respective memories as per the scheme described in section 3.2. The local controller ensures that the Read/ Write signal must go along with the respective addresses. The result of the address generator has been discussed in the next section.

Interconnect Logic: The interconnect logic ensures that there should not be any bus data conflicts i.e. in INPUT data or OUPUT data of the memories. Since these memories either will be in writing or reading mode (out of which two are in reading mode & one in writing mode at a time). The control signal of the memories will change accordingly. The fig. (3.4) shows the logic scheme designed in VHDL. The control signals  $C_{11}$ ,  $C_{12}$  &  $C_{13}$  are the signals used by the memory 1 as chip select ( $\overline{CS}$ ), output enable ( $\overline{OE}$ ) & input/ output (I/ O) respectably. Similarly  $C_{21}$ ,  $C_{22}$  &  $C_{23}$  are the control signals for the memory 2 and  $C_{31}$ ,  $C_{32}$  &  $C_{33}$  are the control signals for the memory 3. These control signals have been generated by using finite state machine (FSM) as per the scheme shown in fig. (3.5). The integration of above two schemes gives the complete design of interconnect logic.

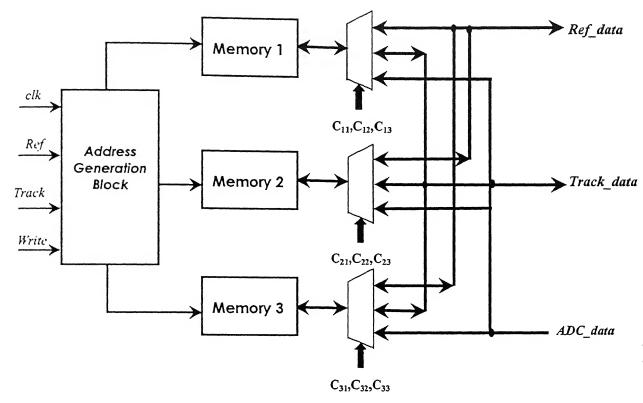


Fig. (3.4): Interconnect Logic for Correlation Tracker

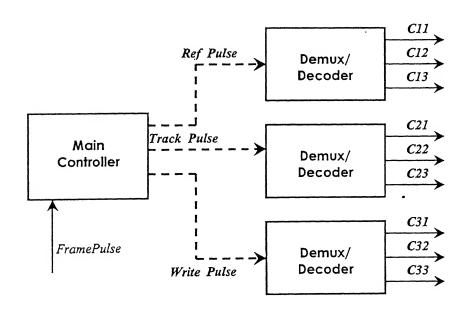


Fig. (3.5): scheme for the implementation of memory Controller

Computational logic: The computation logic uses the sequential similarity detection algorithm (SSDA) for the calculation of the image registration point. It computes the absolute difference of the reference data & track data (or search data) for the  $16 \times 16 = 256$  pixels i.e. the one location. Likewise it will compute the value for all the search position.

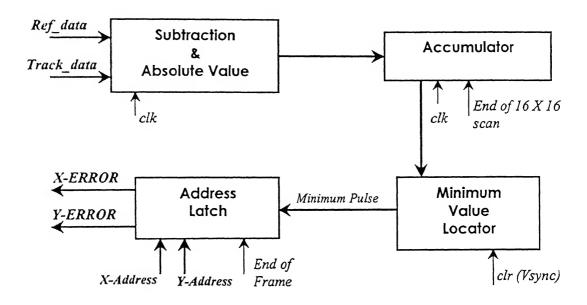


Fig. (3.6): Scheme of Computational Logic for Correlation Tracking

The fig. (3.6) shows the implementation scheme of computational logic. First it will compute the subtracted and absolute value of one pixel of reference data with track data, likewise it will compute for the whole 16 x 16 image size. Then the accumulated value will go to the minimum value locator, which will store the minimum value among the all search area. It will generate the minimum pulse whenever it gets the new minimum value. This minimum pulse uses as a clock to latch the X & Y address at the minimum value. This will be the X-error & Y-error signal.

**3.4 Results and Discussion:** The simulation results of address generator for the correlation tracking implementation are shown in the fig. (3.7). The memory 2 is in the reading mode and gives the reference data from the center of the image size  $64 \times 64$ . So it starts with (1818)<sub>H</sub> and goes up to (1827)<sub>H</sub>. Then it start with (1819)<sub>H</sub>, till it reaches the address (2727)<sub>H</sub> which complete the  $16 \times 16$  block as shown in fig. (3.8).

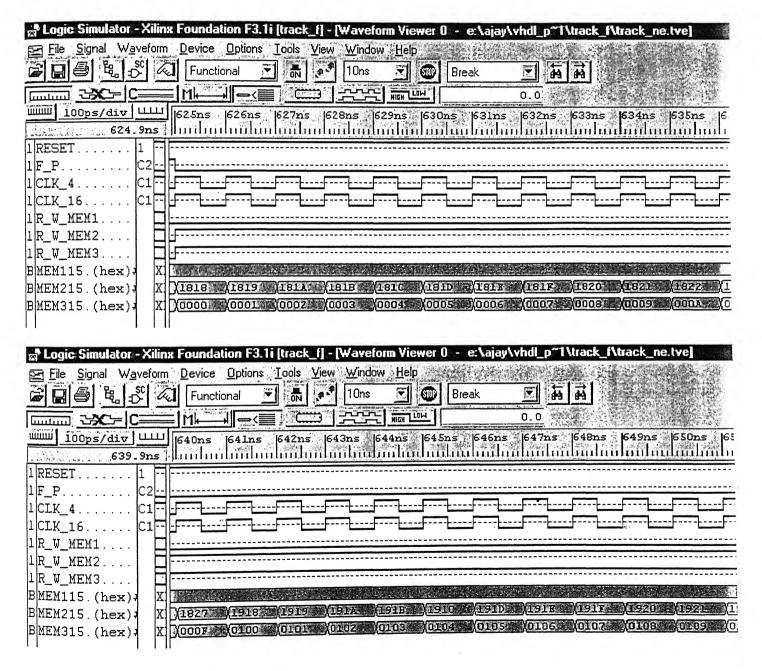


Fig. (3.7): Simulation result of the address Generator

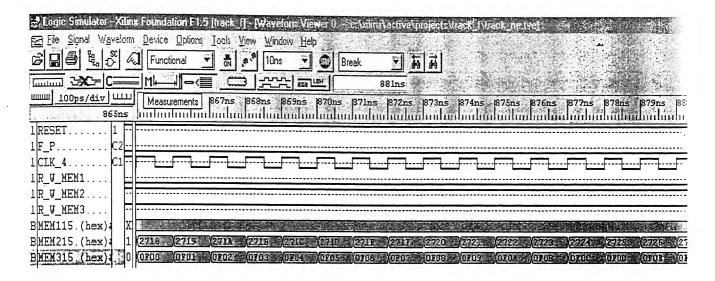


Fig. (3.8): Simulation result of the address Generator shows the completion of 16 x16 block

Fig. (3.9) shows the simulation result of the interconnect logic. This shows the various control signals used by the three memories are as per the detail described above.

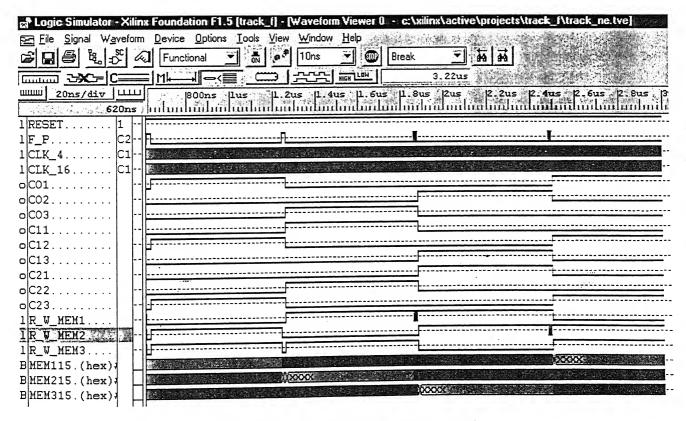


Fig. (3.9): Simulation result of the interconnect logic

This figure also shows the read/ write signals (as R\_W\_MEM1, R\_W\_MEM2 and R\_W\_MEM3) used by the memories. It clearly shows that that the memories change their addresses with change in the frame pulse. While the memories are in the writing mode the address comes with the track gate designed separately. The fig. (3.10) shows that there is no data conflict between the memories INPUT (i.e. ADC data) and OUTPUT data. While the memories accept the data while in the writing mode, otherwise in the high impedance state.

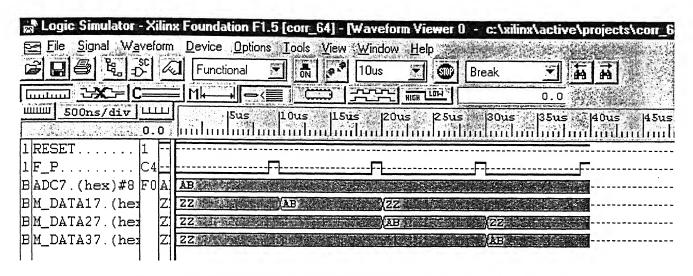


Fig. (3.10): Simulation result of the ADC input data

Logic Simulator - Xilinx Foundation F1.5 [comput] - [Waveform Viewer 0 - c:\xilinx\active\projects\comput\manav.tve]
File Signal Waveform Device Options Iooks View Window Help     ∴
□ □ □ □ □ □ □ □ □ □ □ □ □ □ □ □ □ □ □
THE THE TEST NEW TOWN 12.2865us
100ns/div 1111   pus   2us   3us   4us   5us   6us   7us   8us   9us   10us   11us   12us   13us
i CLK C1
i FRAME   f
BA7(hex)#8 F1 F F0 (F2). AB (FD
BB7(hex)#8 F0 A AX
BSUB7.(hex)#8 5 46 (48) 5 (01) (53)
BACC18.(hex)#; 2 00000 222700 22255¥ 003571 297.54
BMIN18.(hex)# 2 7FFFF (227D0 )2255E (0357A )7FFFF (2975E
I Q 3 3 3 3 3
B ADD12.(hex)#1 0
BEX7(hex)#8 2 E1 26
BEY7(hex)#8 1 71 272
1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1

Fig. (3.11): Simulation result of the computational logic (SSDA method)

Fig. (3.11) shows simulation results for the computational logic. The reference and track data get subtracted & its absolute value gets accumulated for  $16 \times 16$  image size. Whenever the minimum value comes the minimum pulse (i.e. Q) goes high and the corresponding address gets latched. On the occurrence of frame pulse the latch address will be out. The final X-error (shown as EX7) and Y-error (shown as EY7) are taken from the center of the reference image.

The comprehensive synthesis and implementation result of the correlation tracking algorithm into FPGA XC4020E is given below:

Number of CLBs : 760out of 784

CLB Flip Flops : 586

4 input LUTs : 1232 (1 used as route-through)
3 input LUTs : 164 (22 used as route-through)

Number of bonded IOBs : 96 out of 192

IOB Flops :0
IOB Latches :0

Number of TBUFs : 72 out of 1680

Total equivalent gate count for design : 18356 Additional JTAG gate count for IOBs : 4608

The detailed map & pad reports are attached in the appendix-1.

### **CONCLUSION**

In the present work design for an FPGA based implementation of Auto TV tracker has been discussed. Two algorithms centroid and correlation have been considered as tracking techniques. These cater to almost all the scenarios for ground as well as air borne defense systems. For the hardware realization of this system, preliminary design has been done using block diagram approach to ease the comprehension of the design. The complete design has been done using Hardware Description Language (VHDL), so that design modification takes less time.

The centroid algorithm implemented in this work uses the loop-pipelining concept to carry out division. This helps to use one clock (4MHz) for all the process to cater the real-time requirement. It avoids synchronization problem as well. The correlation algorithm implemented in this work uses the three-memory approach, which makes the system capable for use in real-time with the latency of one frame.

# **FUTURE SCOPE**

In the present work the implemented algorithm i.e. centroid and correlation have been designed to cater only for single target. For the coming era, the defense scenario may be there to cater for multiple targets. So these designed algorithm could be modified for the multiple target image tracking. Centroid algorithms designed for the present work can be further modified by using sharing of resources through scheduling scheme. However this will have area-delay tradeoff. In this work the memories (SRAM) used for the implementation of correlation technique are external. Now a days since the embedded FPGA are coming whose core are design in such a way that it can be programmed for the design which uses the large memory block, so the external memory can be removed. This concept will help in the reduction of circuit size and PCB interconnects, which will make the system fast as well as compact. If the communication interface (e.g. RS-422) is incorporated in the Auto TV tracker circuit then it could be interface with any standard system, as these days most of the systems are coming with this kind of interface.

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# APPENDIX - I

# FPGA Implementation Report of Centroid Tracking Algorithm:

#### MAP Report:

```
Xilinx Mapping Report File for Design "trakcen1"
Copyright (c) 1995-1998 Xilinx, Inc. All rights reserved.
```

```
Design Information
Command Line : map -p xc4008e-3-pc84 -o map.ncd trakcen1.ngd
trakcen1.pcf
Target Device : x4008e
Target Package : pc84
Target Speed : -3
Mapper Version: xc4000e -- M1.5.19
Mapped Date : Thu Nov 08 16:15:44 2001
Design Summary
_____
  Number of errors:
                         Ω
  Number of warnings: 1
Number of CLBs: 233 out of 324 72%
CLB Flip Flops: 190
     4 input LUTs: 171
3 input LUTs: 54 (15 used as route-throughs)
   Number of bonded IOBs: 25 out of 61 40%
     IOB Latches:
     IOB Flops:
                         0
  Number of clock IOB pads: 8 out of 8 100%
                                        4 100%
  Number of primary CLKs:
                            4 out of
  Number of secondary CLKs: 4 out of
                                        4 100%
                            48 out of 720
  Number of TBUFs:
                                             6ક
  Number of RPM macros:
Total equivalent gate count for design: 6425
Additional JTAG gate count for IOBs: 1200
Table of Contents
______
Section 1 - Errors
```

```
Section 2 - Warnings
Section 3 - Design Attributes
Section 4 - Removed Logic Summary
Section 5 - Removed Logic
Section 6 - Added Logic
Section 7 - Expanded Logic
Section 8 - Signal Cross-Reference
Section 9 - Symbol Cross-Reference
Section 10 - IOB Properties
```

Section 11 - RPMs

Section 12 - Guide Report

```
Section 1 - Errors
______
Section 2 - Warnings
_______
Section 3 - Design Attributes
______
Section 4 - Removed Logic Summary
______
 15 block(s) optimized away
Section 5 - Removed Logic
_____
Optimized Block(s):
TYPE
         BLOCK
GND
         H1/I_from_net_GND
GND
         I_from_net_GND
         U12/templ_reg<15>
DFF
DFF
         U7/q_reg<7>
X ZERO
              U7/N67.ZERO
GND
         H1/U1/N12_GND_0
X ZERO
              U12/N8.ZERO
         U12/temp1_reg<10>
DFF
         U12/temp1 reg<l1>
DFF
         U12/temp1 reg<12>
DFF
         U12/templ_reg<13>
DFF
         U12/temp1_reg<14>
DFF
         U12/temp1 reg<8>
DFF
DFF
         U12/templ reg<9>
              U9/N10.ZERO
X ZERO
Section 6 - Added Logic
----
Section 7 - Expanded Logic
______
To enable this section, set the environment variable MAP_REPORT_DETAIL
to TRUE
and rerun MAP.
Section 8 - Signal Cross-Reference
To enable this section, set the environment variable MAP_REPORT_DETAIL
to TRUE
and rerun MAP.
Section 9 - Symbol Cross-Reference
------
To enable this section, set the environment variable MAP_REPORT_DETAIL
to TRUE
and rerun MAP.
```

Section 10 - IOB Properties

```
_______
X ADD<0> (IOB) : SLEW=SLOW
X_ADD<1> (IOB) : SLEW=SLOW
X ADD<2> (IOB) : SLEW=SLOW
X ADD<3> (IOB) : SLEW=SLOW
X_ADD<4> (IOB) : SLEW=SLOW
X ADD<5> (IOB) : SLEW=SLOW
X ADD<6> (IOB) : SLEW=SLOW
X_ADD<7> (IOB) : SLEW=SLOW
Y ADD<0> (IOB) : SLEW=SLOW
Y ADD<1> (IOB) : SLEW=SLOW
Y_ADD<2> (IOB) : SLEW=SLOW
Y ADD<3> (IOB) : SLEW=SLOW
Y ADD<4> (IOB) : SLEW=SLOW
Y_ADD<5> (IOB) : SLEW=SLOW
Y ADD<6> (IOB) : SLEW=SLOW
Y ADD<7> (IOB) : SLEW=SLOW
Section 11 - RPMs
H1/U1/SUB_RPL_16_4_0
                                   - 9 comps
H1/U3/INC RPL 8 3 0
                                    - 5 comps
H1/U3/DEC RPL 8 4 1
                                    - 5 comps
U11/INC_RPL_8_3_0
                                    - 5 comps
                                    - 9 comps
U12/add rpl 16 2 0
U5/inc_rpl_16_3_0
                                   - 9 comps
U7/inc_rpl_8_3_0
                                   - 5 comps
                                   - 5 comps
U8/inc rpl 8 3 0
U9/sub_rpl_8_1_0
                                    - 5 comps
Section 12 - Guide Report
Guide not run on this design.
Pad Report:
PAR: Xilinx Place And Route M1.5.19.
Copyright (c) 1995-1998 Xilinx, Inc. All rights reserved.
Thu Nov 08 16:16:07 2001
Xilinx PAD Specification File
******
Input file:
               map.ncd
Output file:
               trakcen1.ncd
               xc4008e
Part type:
Speed grade: -3
Package:
                pc84
Thu Nov 08 16:16:07 2001
Pinout by Pin Number:
+-----
```

Pin Numb	per   Pin Name	Direction
+   P1	GND	
P2	VCC	
P3	VCC	
	1	
P4	1	
P5	1	
P6		
P7		
P8		
P9		
P10	TVP	INPUT
P11	VCC	1
P12	GND	1
P13	TVPVER	INPUT
P14	X_ADD<7>	OUTPUT
P15	X_ADD<1>	OUTPUT
P16	X_ADD<4>	OUTPUT
P17	Y_ADD<7>	OUTPUT
P18	Y_ADD<6>	OUTPUT
P19	Y_ADD<0>	OUTPUT
P20	Y_ADD<3>	OUTPUT
P21	GND	
P22	VCC	
P23	Y_ADD<4>	OUTPUT
P24	X_ADD<3>	OUTPUT
P25	1	
P26	1	
1		

P27	1	1
P28	1	1
   P29	VSYNC	INPUT
P30	Ml	1
P31	GND	1
P32	MO	1
P33	VCC	1
P34	M2	1 1
P35	CLK_4	INPUT
P36		1
P37	1	1
P38	SEL	INPUT
P39		1
P40		1
P41		1
P42	VCC	1
P43	GND	1 1
P44		1
P45		1 1
P46		1 1
P47		1
P48		1 1
P49		1 . 1
   P50		1
   P51 	HSYNC	INPUT
P52	GND	1
P53	DONE	1
P54	VCC	1

	P55	/PROG	1		
	P56		ļ		1
	P57	RESET	ļ	INPUT	1
	P58		1		
	P59		1		]
	P60		ļ		
	P61	X_ADD<2>	ļ	OUTPUT	
	P62		1		
and distances and	P63	VCC			
	P64	GND			
	P65	Y_ADD<2>		OUTPUT	
The second second	P66	Y_ADD<1>		OUTPUT	
-	P67	Y_ADD<5>		OUTPUT	
	P68	X_ADD<0>		OUTPUT	
-	P69	X_ADD<5>		OUTPUT	
-	P70	X_ADD<6>		OUTPUT	
and the same of	P71	1	1		
and decrees the	P72	OEDIVIN		INPUT	
-	P73	CCTK			
The second second	P74	VCC			
-	P75	TDO			
	P76	GND			
	P77	1			
	P78	win		INPUT	
	P79	1 .	1		1 .
	P80	1			1
	P81	1	1		]
	P82	1	l		
ı					

```
P83
P84
TDO
# Pinout constraints listing
# These constraints are in PCF grammar format
# and may be cut and pasted into the PCF file
# after the "SCHEMATIC END ; " statement to
# preserve this pinout for future design iterations.
COMP "CLK_4" LOCATE = SITE "P35" ;
COMP "HSYNC" LOCATE = SITE "P51" ;
COMP "OEDIVIN" LOCATE = SITE "P72" ;
COMP "RESET" LOCATE = SITE "P57" ;
COMP "SEL" LOCATE = SITE "P38" ;
COMP "TVP" LOCATE = SITE "P10" ;
COMP "TVPVER" LOCATE = SITE "P13" ;
COMP "VSYNC" LOCATE = SITE "P29" ;
COMP "WIN" LOCATE = SITE "P78" ;
COMP "X ADD<0>" LOCATE = SITE "P68" ;
COMP "X ADD<1>" LOCATE = SITE "P15"
COMP "X ADD<2>" LOCATE = SITE "P61"
COMP "X ADD<3>" LOCATE = SITE "P24"
COMP "X_ADD<4>" LOCATE = SITE "P16"
COMP "X ADD<5>" LOCATE = SITE "P69" ;
COMP "X ADD<6>" LOCATE = SITE "P70"
COMP "X_ADD<7>" LOCATE = SITE "P14"
COMP "Y_ADD<0>" LOCATE = SITE "P19"
COMP "Y ADD<1>" LOCATE = SITE "P66" ;
COMP "Y ADD<2>" LOCATE = SITE "P65"
COMP "Y_ADD<3>" LOCATE = SITE "P20" ;
COMP "Y ADD<4>" LOCATE = SITE "P23" ;
COMP "Y_ADD<5>" LOCATE = SITE "P67"
COMP "Y ADD<6>" LOCATE = SITE "P18"
COMP "Y ADD<7>" LOCATE = SITE "P17" ;
#
```

# FPGA Implementation Report of Correlation Tracking Algorithm:

#### MAP Report:

```
Xilinx Mapping Report File for Design 'alh corr'
Copyright (c) 1995-2000 Xilinx, Inc. All rights reserved.
Design Information
______
Command Line : mlmap -p xc4020e-1-pg223 -o map.ncd alh_corr.ngd
alh corr.pcf
Target Device : x4020e
Target Package: pg223
Target Speed : -1
Mapper Version : xc4000e -- D.19
Design Summary
Number of errors:
  Number of CLBs:
CLB Flin D'
                           760 out of 784 96%
                       586
     CLB Flip Flops:
     4 input LUTs: 1232 (1 used as route-throughs)
3 input LUTs: 164 (22 used as route-throughs)
   Number of bonded IOBs: 96 out of 192 50%
      IOB Flops:
     IOB Latches:
                         0
  Number of TBUFs:
                             72 out of 1680
                                              4 %
   15 unrelated functions packed into 15 CLBs.
   (1% of the CLBs used are affected.)
Total equivalent gate count for design: 18356
Additional JTAG gate count for IOBs: 4608
Table of Contents
Section 1 - Errors
Section 2 - Warnings
Section 3 - Design Attributes
Section 4 - Removed Logic Summary
Section 5 - Removed Logic
Section 6 - Added Logic
Section 7 - Expanded Logic
Section 8 - Signal Cross-Reference
Section 9 - Symbol Cross-Reference
Section 10 - IOB Properties
Section 11 - RPMs
Section 12 - Guide Report
Section 13 - Area Group Summary
Section 1 - Errors
------
Section 2 - Warnings
______
Section 3 - Design Attributes
-------
Section 4 - Removed Logic Summary
242 block(s) removed
 232 block(s) optimized away
```

```
Section 5 - Removed Logic
The trimmed logic reported below is either:
  1. part of a cycle
  2. part of disabled logic
  3. a side-effect of other trimmed logic
Unused block "U21/C954" (GND) removed.
 Unused block "U25/C80" (GND) removed.
Optimized Block(s):
TYPE
          BLOCK
VCC
          $150
VCC
           $174
INV
          U17/C257
          U20/C74
INV
INV
          U22/C132
INV
          U23/C126
X ZERO
                 H1/U1/N194.ZERO
X_ONE
                H1/U1/_Logic1 .ONE
X ZERO
                H1/U2/N54.ZERO
X ZERO
                H2/U1/N194.ZERO
X ONE
                H2/U1/_Logic1 .ONE
X ZERO
                H2/U2/N54.ZERO
X ZERO
                H3/U1/N194.ZERO
X ONE
                H3/U1/ Logic1 .ONE
X ZERO
                H3/U2/N54.ZERO
          U15/C217
GND
GND
          U16/C167
VCC
          U17/C308
GND
           U17/C309
VCC
           U21/C912
To enable printing of redundant blocks removed and signals merged, set
the
detailed map report option and rerun map.
Section 6 - Added Logic
Section 7 - Expanded Logic
-----
To enable this section, set the detailed map report option and rerun
Section 8 - Signal Cross-Reference
To enable this section, set the detailed map report option and rerun
map.
Section 9 - Symbol Cross-Reference
To enable this section, set the detailed map report option and rerun
map.
Section 10 - IOB Properties
"CNT 254" (IOB) : SLEW=SLOW
"E_X<0>" (IOB) : SLEW=SLOW
"E X<1>" (IOB) : SLEW=SLOW
                                            पुरुषोत्तम का गीमान केलकर पुस्तकालक
"E X<2>" (IOB) : SLEW=SLOW
                                            भारतीय प्रौद्योति शिल्लां कानपुर
अवाष्ति क्रः A 13.792.5
"E X<3>" (IOB) : SLEW=SLOW
"E X<4>" (IOB) : SLEW=SLOW
```

242 signal(s) removed

```
"E X<5>" (IOB) : SLEW=SLOW
"E X<6>" (IOB) : SLEW=SLOW
"E X<7>" (IOB) : SLEW=SLOW
"E Y<0>" (ICB) : SLEW=SLOW
"E Y<1>" (IOB) : SLEW=SLOW
"E Y<2>" (IOB) : SLEW=SLOW
"E Y<3>" (IOB) : SLEW=SLOW
"E Y<4>" (IOB) : SLEW=SLOW
"E Y<5>" (IOB) : SLEW=SLOW
"E Y<6>" (IOB) : SLEW=SLOW
"E Y<7>" (IOB) : SLEW=SLOW
"MEM10" (IOB) : SLEW=SLOW
"MEM11" (IOB) : SLEW=SLOW
"MEM110" (IOB) : SLEW=SLOW
"MEM111" (IOB) : SLEW=SLOW
"MEM12" (IOB) : SLEW=SLOW
"MEM13" (IOB) : SLEW=SLOW
"MEM14" (IOB) : SLEW=SLOW
"MEM15" (IOB) : SLEW=SLOW
"MEM16" (IOB) : SLEW=SLOW
"MEM17" (IOB) : SLEW=SLOW
"MEM18" (IOB) : SLEW=SLOW
"MEM19" (IOB) : SLEW=SLOW
"MEM20" (IOB) : SLEW=SLOW
"MEM21" (IOB) : SLEW=SLOW
"MEM210" (IOB) : SLEW=SLOW
"MEM211" (IOB) : SLEW=SLOW
"MEM22" (IOB) : SLEW=SLOW
"MEM23" (IOB) : SLEW=SLOW
"MEM24" (IOB) : SLEW=SLOW
"MEM25" (IOB) : SLEW=SLOW
"MEM26" (IOB) : SLEW=SLOW
"MEM27" (IOB) : SLEW=SLOW
"MEM28" (IOB) : SLEW=SLOW
"MEM29" (IOB) : SLEW=SLOW
"MEM30" (IOB) : SLEW=SLOW
"MEM31" (IOB) : SLEW=SLOW
"MEM310" (IOB) : SLEW=SLOW
"MEM311" (IOB) : SLEW=SLOW
"MEM32" (IOB) : SLEW=SLOW
"MEM33" (IOB) : SLEW=SLOW
"MEM34" (IOB) : SLEW=SLOW
"MEM35" (IOB) : SLEW=SLOW
"MEM36" (IOB) : SLEW=SLOW
"MEM37" (IOB) : SLEW=SLOW
"MEM38" (IOB) : SLEW=SLOW
"MEM39" (IOB) : SLEW=SLOW
"M DATA<10>" (IOB) : SLEW=SLOW
"M DATA<11>" (IOB) : SLEW=SLOW
"M DATA<12>" (IOB) : SLEW=SLOW
"M DATA<13>" (IOB) : SLEW=SLOW
"M DATA<14>" (IOB) : SLEW=SLOW
"M DATA<15>" (IOB) : SLEW=SLOW
"M DATA<16>" (IOB) : SLEW=SLOW
"M_DATA<17>" (IOB) : SLEW=SLOW
"M_DATA<20>" (IOB) : SLEW=SLOW
"M_DATA<21>" (IOB) : SLEW=SLOW
```

```
"M_DATA<22>" (IOB) : SLEW=SLOW
"M DATA<23>" (IOB) : SLEW=SLOW
"M DATA<24>" (IOB) : SLEW=SLOW
"M DATA<25>" (IOB) : SLEW=SLOW
"M DATA<26>" (IOB) : SLEW=SLOW
"M_DATA<27>" (IOB) : SLEW=SLOW
"M_DATA<30>" (IOB) : SLEW=SLOW
"M DATA<31>" (IOB) : SLEW=SLOW
"M_DATA<32>" (IOB) : SLEW=SLOW
"M DATA<33>" (IOB) : SLEW=SLOW
"M DATA<34>" (IOB) : SLEW=SLOW
"M DATA<35>" (IOB) : SLEW=SLOW
"M_DATA<36>" (IOB) : SLEW=SLOW
"M DATA<37>" (IOB) : SLEW=SLOW
"R_W_MEM1" (IOB) : SLEW=SLOW
"R_W_MEM2" (IOB) : SLEW=SLOW
"R W MEM3" (IOB) : SLEW=SLOW
"THREE" (IOB) : SLEW=SLOW
"TX_DATA" (IOB) : SLEW=SLOW
"WIN 64" (IOB) : SLEW=SLOW
Section 11 - RPMs
Section 12 - Guide Report
-----
Guide not run on this design.
```

#### Pad Report:

Ŧ		•		•		
100	+ Pin Number onstraint	and the same of th	Pin Name	  -	Direction	 
+	+	-+		т :	•	
	A2		(TDI)			
	А3	1	MEM35	l	OUTPUT	!
	A4	l	MEM27		OUTPUT	l
	A5	1	E_Y<4>	İ	OUTPUT	l
	A6	1	MEM311		OUTPUT	

A7	MEM22	1	OUTPUT	
A8	MEM14	1	OUTPUT	1
A9	MEM21	1	OUTPUT	
A10	MEM210	1	OUTPUT	1
All	MEM24	1	OUTPUT	1
A12	E_X<5>	1	OUTPUT	
A13	RESET		INPUT	1
A14		1	UNUSED	
A15		1	UNUSED	ŀ
A16		1	UNUSED	
A17			UNUSED	
A18	(MO)			1
Bl		-	UNUSED	
B2		1	UNUSED	
B3			UNUSED	
B4	(TCK)	1		
B5	MEM34	1	OUTPUT	1
B6	E_Y<1>	1	OUTPUT	
B7	E_Y<3> (TMS)	1	OUTPUT	
B8	MEM16		OUTPUT	
B9	MEM26	}	OUTPUT	
B10	MEM110		OUTPUT	1
B11	MEM3 3		OUTPUT	1
B12	E_X<6>	-	OUTPUT	1
B13		1	UNUSED	
B14		ļ	UNUSED	
B15			UNUSED	
B16	CTK	1	INPUT	
l				

B17		UNUSED
B18		UNUSED
Cl	MEM18	OUTPUT
C2		UNUSED
C3		UNUSED
C4		UNUSED
C5		UNUSED
C6	MEM37	OUTPUT
C7	(GND)	1
C8	MEM15	OUTPUT
C9	MEM11	OUTPUT
C10	MEM310	OUTPUT
C11	MEM13	OUTPUT
C12	(GND)	!
C13		UNUSED
C14		UNUSED
C15	(M1)	
C16	(M2)	1
C17		UNUSED
C18		UNUSED
D1	MEM36	OUTPUT
D2	MEM17	OUTPUT
D3	(VCC)	
D4	(GND)	1
D5	MEM211	OUTPUT
D6	MEM31	OUTPUT
ם 7	MEM111	OUTPUT
D8	MEM20	OUTPUT
1		

D9	(GND)	1	
D10	(VCC)	1	
D11	MEM23	OUTPUT	1
D12	MEM25	OUTPUT	
D13	MEM10	OUTPUT	1
D14		UNUSED	1
D15	(GND)	1	1
D16	(VCC)	1	1
D17		UNUSED	
D18		UNUSED	
E1	MEM39	OUTPUT	
E2	MEM28	OUTPUT	ı
E3		UNUSED	.
E4		UNUSED	ı
E15		UNUSED	1
   E16		UNUSED	
E17		UNUSED	1
E18		UNUSED	1
F1	MEM29	OUTPUT	
   F2	E_Y<0>	OUTPUT	[
F3	MEM38	OUTPUT	1
F4	MEM30	OUTPUT	1
F15		UNUSED	1
F16		UNUSED	1
F17		UNUSED	1
F18	F_P	INPUT	1
G1	M_DATA<20>	BIDIR	1
G2	MEM32	OUTPUT	

G3	(GND)		
G4	E_Y<2>	OUTPUT	1
G15	R_W_MEM3	OUTPUT	
G16	(GND)		
G17		UNUSED	
G18		UNUSED	
H1	MEM12	OUTPUT	
H2	M_DATA<22>	BIDIR	I
Н3	R_W_MEM2	OUTPUT	1
H4	MEM19	OUTPUT	1
H15	M_DATA<31>	BIDIR	1
H16		UNUSED	
H17	ADC<7>	INPUT	
H18	M_DATA<37>	BIDIR	
J1	M_DATA<34>	BIDIR	1
J2	M_DATA<36>	BIDIR	1
J3	M_DATA<30>	BIDIR	
J4	(VCC)	!	
J15	(VCC)		1
   J16	M_DATA<35>	BIDIR	1
J17	ADC<5>	INPUT	
   J18	M_DATA<33>	BIDIR	
K1	M_DATA<24>	BIDIR	
K2	M_DATA<26>	BIDIR	
K3	M_DATA<32>	BIDIR	
   K4	(GND)	I	
   K15	(GND)	1	-
   K16	M_DATA<25>	BIDIR	

K17	M_DATA<23>	BIDIR
К18	ADC<1>	INPUT
L1	ADC<4>	INPUT
L2	E_X<0>	OUTPUT
L3	ADC<2>	INPUT
L4		UNUSED
L15	THREE	OUTPUT
L16	ADC<3>	INPUT
L17	M_DATA<27>	BIDIR
L18	M_DATA<21>	BIDIR
MI	ADC<0>	INPUT
M2	ADC<6>	INPUT
МЗ	(GND)	1
M4		UNUSED
M15	M_DATA<15>	BIDIR
M16	(GND)	1
M17	ļ	UNUSED
M18	M_DATA<17>	BIDIR
Nl		UNUSED
N2		UNUSED
из		UNUSED
N4		UNUSED
N15	CNT_254	OUTPUT
N16		UNUSED
N17		UNUSED
N18		UNUSED
		1
Pl		UNUSED

	P3			UNUSED	
	P4		[	UNUSED	
	   P15	R_W_MEM1	1	OUTPUT	
	P16			UNUSED	
	P17		1	UNUSED	I
	P18		1	UNUSED	1
	R1			UNUSED	
	R2			UNUSED	1
	R3	(GND)	1		
	R4	(VCC)			
	R5		-	UNUSED	
	R6	M_DATA<12>		BIDIR	
	R7			UNUSED	
-	R8	M_DATA<16>	1	BIDIR	
-	R9	(GND)			
	R10	(VCC)	l		
	R11			UNUSED	
	R12		l	UNUSED	
	R13		1	UNUSED	1
-	R14		ļ	UNUSED	١
-	R15	(VCC)			
	R16	(GND)			1
	R17			UNUSED	1
-	R18			UNUSED	1
	Tl		1	UNUSED	1
	T2		1	UNUSED	1
	тз			UNUSED	
	T4			UNUSED	
- 1					

	T5		UNUSED	
	T6		UNUSED	
-	<b>T</b> 7	(GND)	1	1
	T8	E_X<2>	OUTPUT	1
	Т9	M_DATA<10>	BIDIR	1
	T10	E_X<3>	OUTPUT	1
	Tll	WIN_64	OUTPUT	
	T12	(GND)		
	T13		UNUSED	1
	T14		UNUSED	1
	T15		UNUSED	1
	T16		UNUSED	
	T17		UNUSED	
	T18		UNUSED	
	ן עו		UNUSED	1
	   U2	(TDO)		1
	U3		UNUSED	1
	U4		UNUSED	
	   U5		UNUSED	1
	   U6		UNUSED	1
	ן עד	E_Y<6>	OUTPUT	
	U8	E_X<1>	OUTPUT	1
	   U9	E_X<7>	OUTPUT	
	U10		UNUSED	1
	   Ull		UNUSED	
	   U12		UNUSED	
	   U13		UNUSED	1
	U14	·	UNUSED	1
	1			

```
U15
           | ---
                                        UNUSED
U16
                                        UNUSED
U17
            (DONE)
Uls
                                        UNUSED
Vl
           (CCLK)
V2
                                       UNUSED
V3
                                       UNUSED
V4
                                       UNUSED
V5
           M DATA<14>
                                       BIDIR
V6
           E Y<5>
                                       OUTPUT
V7
           E Y<7>
                                        OUTPUT
V8
           E X<4>
                                       OUTPUT
V9
           TRACK ENG
                                        INPUT
Vlo
           HSY
                                        INPUT
Vll
           ---
                                        UNUSED
V12
           M DATA<11>
                                       BIDIR
V13
           | TX DATA
                                       OUTPUT
V14
           M DATA<13>
                                        BIDIR
V15
                                        UNUSED
V16
                                        UNUSED
V17
                                        UNUSED
V18
           (/PROG)
```

```
#
# To preserve the pinout above for future design iterations,
# simply run "Lock Pins..." from the Design Manager's Design
# menu, or invoke PIN2UCF from the command line. The location
constraints
# above will be written into your specified UCF file. (The constraints
# listed below are in PCF format and cannot be directly used in the UCF
file).
#
COMP "ADC<0>" LOCATE = SITE "M1";
COMP "ADC<1>" LOCATE = SITE "K18";
```

```
COMP "ADC<2>" LOCATE = SITE "L3" ;
COMP "ADC<3>" LOCATE = SITE "L16" ;
COMP "ADC<4>" LOCATE = SITE "L1" ;
COMP "ADC<5>" LOCATE = SITE "J17" ;
COMP "ADC<6>" LOCATE = SITE "M2" ;
COMP "ADC<7>" LOCATE = SITE "H17" ;
COMP "CLK" LOCATE = SITE "B16" ;
COMP "CNT 254" LOCATE = SITE "N15" ;
COMP "E X<0>" LOCATE = SITE "L2" ;
COMP "E X<1>" LOCATE = SITE "U8" ;
COMP "E X<2>" LOCATE = SITE "T8" ;
COMP "E X<3>" LOCATE = SITE "T10" ;
COMP "E X<4>" LOCATE = SITE "V8" ;
COMP "E X<5>" LOCATE = SITE "A12" ;
COMP "E X<6>" LOCATE = SITE "B12" ;
COMP "E X<7>" LOCATE = SITE "U9" ;
COMP "E Y<0>" LOCATE = SITE "F2" ;
COMP "E Y<1>" LOCATE = SITE "B6" ;
COMP "E Y<2>" LOCATE = SITE "G4" ;
COMP "E Y<3>" LOCATE = SITE "B7" ;
COMP "E Y<4>" LOCATE = SITE "A5" ;
COMP "E Y<5>" LOCATE = SITE "V6" ;
COMP "E Y<6>" LOCATE = SITE "U7" ;
COMP "E Y<7>" LOCATE = SITE "V7" ;
COMP "F P" LOCATE = SITE "F18" ;
COMP "HSY" LOCATE = SITE "V10" ;
COMP "MEM10" LOCATE = SITE "D13" ;
COMP "MEM11" LOCATE = SITE "C9" ;
COMP "MEM110" LOCATE = SITE "B10" ;
COMP "MEM111" LOCATE = SITE "D7" ;
COMP "MEM12" LOCATE = SITE "H1" ;
COMP "MEM13" LOCATE = SITE "C11" ;
COMP "MEM14" LOCATE = SITE "A8" ;
COMP "MEM15" LOCATE = SITE "C8" ;
COMP "MEM16" LOCATE = SITE "B8" ;
COMP "MEM17" LOCATE = SITE "D2" ;
COMP "MEM18" LOCATE = SITE "C1" ;
COMP "MEM19" LOCATE = SITE "H4" ;
COMP "MEM20" LOCATE = SITE "D8" ;
COMP "MEM21" LOCATE = SITE "A9" ;
COMP "MEM210" LOCATE = SITE "A10" ;
COMP "MEM211" LOCATE = SITE "D5" ;
COMP "MEM22" LOCATE = SITE "A7" ;
COMP "MEM23" LOCATE = SITE "D11" ;
COMP "MEM24" LOCATE = SITE "A11" ;
COMP "MEM25" LOCATE = SITE "D12" ;
COMP "MEM26" LOCATE = SITE "B9" ;
COMP "MEM27" LOCATE = SITE "A4" ;
COMP "MEM28" LOCATE = SITE "E2" ;
COMP "MEM29" LOCATE = SITE "F1" ;
COMP "MEM30" LOCATE = SITE "F4" ;
COMP "MEM31" LOCATE = SITE "D6" ;
COMP "MEM310" LOCATE = SITE "C10" ;
COMP "MEM311" LOCATE = SITE "A6" ;
COMP "MEM32" LOCATE = SITE "G2" ;
COMP "MEM33" LOCATE = SITE "B11" ;
COMP "MEM34" LOCATE = SITE "B5" ;
```

```
COMP "MEM35" LOCATE = SITE "A3" ;
COMP "MEM36" LOCATE = SITE "D1" :
COMP "MEM37" LOCATE = SITE "C6" ;
COMP "MEM38" LOCATE = SITE "F3"
COMP "MEM39" LOCATE = SITE "E1" ;
COMP "M_DATA<10>" LOCATE = SITE "T9" ;
COMP "M DATA<11>" LOCATE = SITE "V12" ;
COMP "M_DATA<12>" LOCATE = SITE "R6" ;
COMP "M_DATA<13>" LOCATE = SITE "V14" ;
COMP "M_DATA<14>" LOCATE = SITE "V5" ;
COMP "M DATA<15>" LOCATE = SITE "M15" ;
COMP "M DATA<16>" LOCATE = SITE "R8" ;
COMP "M DATA<17>" LOCATE = SITE "M18" ;
COMP "M DATA<20>" LOCATE = SITE "G1" ;
COMP "M DATA<21>" LOCATE = SITE "L18" ;
COMP "M DATA<22>" LOCATE = SITE "H2" ;
COMP "M_DATA<23>" LOCATE = SITE "K17" ;
COMP "M DATA<24>" LOCATE = SITE "K1" ;
COMP "M DATA<25>" LOCATE = SITE "K16" ;
COMP "M DATA<26>" LOCATE = SITE "K2" ;
COMP "M DATA<27>" LOCATE = SITE "L17" ;
COMP "M DATA<30>" LOCATE = SITE "J3" ;
COMP "M_DATA<31>" LOCATE = SITE "H15" ;
COMP "M DATA<32>" LOCATE = SITE "K3" ;
COMP "M DATA<33>" LOCATE = SITE "J18" ;
COMP "M DATA<34>" LOCATE = SITE "J1" ;
COMP "M DATA<35>" LOCATE = SITE "J16" ;
COMP "M DATA<36>" LOCATE = SITE "J2" ;
COMP "M DATA<37>" LOCATE = SITE "H18" ;
COMP "RESET" LOCATE = SITE "A13" ;
COMP "R_W_MEM1" LOCATE = SITE "P15" ;
COMP "R W MEM2" LOCATE = SITE "H3" ;
COMP "R W MEM3" LOCATE = SITE "G15" ;
COMP "THREE" LOCATE = SITE "L15" ;
COMP "TRACK ENG" LOCATE = SITE "V9" ;
COMP "TX_DATA" LOCATE = SITE "V13" ;
COMP "WIN 64" LOCATE = SITE "T11" ;
```

# **APPENDIX-II**

#### Brief Description of Xilinx FPGA:

XC4000 Series devices are implemented with a regular, flexible, programmable architecture of Configurable Logic Blocks (CLBs), interconnected by a powerful hierarchy of versatile routing resources, and surrounded by a perimeter of programmable Input/Output Blocks (IOBs). They have generous routing resources to accommodate the most complex interconnect patterns. The devices are customized by loading configuration data into internal memory cells. The FPGA can either actively read its configuration data from an external serial or byte-parallel PROM (master modes), or the configuration data can be written into the FPGA from an external device (slave and peripheral modes). XC4000 Series FPGAs are supported by powerful and sophisticated software, covering every aspect of design from schematic or behavioral entry, floor planning, simulation, automatic block placement and routing of interconnects, to the creation, downloading, and readback of the configuration bit stream. Because Xilinx FPGAs can be reprogrammed an unlimited number of times, they can be used in innovative designs where hardware is changed dynamically, or where hardware must be adapted to different user applications. FPGAs are ideal for shortening design and development cycles, and also offer a cost-effective solution for production rates well beyond 5,000 systems per month. For lowest high-volume unit cost, a design can first be implemented in the XC4000E or XC4000X, then migrated to one of Xilinx' compatible Hardwired maskprogrammed devices.

Device	Logic Cell	Max. Logic Gate	CLB Matrix	Total CLB	No. of Flip/ Flop	Max. us I/ O	
4008E	770	8000	18 x 18	324	936	144	
4020E	1862	20000	28 x 28	784	2016	224	

### Detailed Functional Description of Xilinx FPGA Architecture:

XC4000 Series devices achieve high speed through advanced semiconductor technology and improved architecture. The XC4000E and XC4000X support system clock rates of up to 80 MHz and internal performance in excess of 150 MHz. Compared to older Xilinx FPGA families, XC4000 Series devices are more powerful. They offer on-chip edge-triggered and dual-port RAM, clock enables on I/O flip-flops, and wide-input decoders. They are more versatile in many applications, especially those involving RAM. Design cycles are faster due to a combination of increased routing resources and more sophisticated software.

#### Basic Building Blocks

Xilinx user-programmable gate arrays include two major configurable elements, configurable logic blocks (CLBs) and input/output blocks (IOBs).

- CLBs provide the functional elements for constructing the user's logic.
- IOBs provide the interface between the package pins and internal signal lines.

Three other types of circuits are also available:

- 3-State buffers (TBUFs) driving horizontal longlines are associated with each CLB.
- Wide edge decoders are available around the periphery of each device.
- An on-chip oscillator is provided.

Programmable interconnect resources provide routing paths to connect the inputs and outputs of these configurable elements to the appropriate networks. The functionality of each circuit block is customized during configuration by programming internal static memory cells. The values stored in these memory cells determine the logic functions and interconnections implemented in the FPGA. Each of these available circuits is described in this section.

# Configurable Logic Blocks (CLBs):

Configurable Logic Blocks implement most of the logic in an FPGA. The principal CLB elements are shown in Figure 1. Two 4-input function generators (F and G) offer unrestricted versatility. Most combinatorial logic functions need four

or fewer inputs. However, a third function generator (H) is provided. The H function generator has three inputs. Either zero, one, or two of these inputs can be the outputs of F and G; the other input(s) are from outside the CLB. The CLB can, therefore, implement certain functions of up to nine variables, like parity check or expandable identity comparison of two sets of four inputs. Each CLB contains two storage elements that can be used to store the function generator outputs.

However, the storage elements and function generators can also be used independently. These storage elements can be configured as flip-flops in both XC4000E and XC4000X devices; in the XC4000X they can optionally be configured as latches. DIN can be used as a direct input to either of the two storage elements. H1 can drive the other through the H function generator. Function generator outputs can also drive two outputs independent of the storage element outputs. This versatility increases logic capacity and simplifies routing. Thirteen CLB inputs and four CLB outputs provide access to the function generators and storage elements. These inputs and outputs connect to the programmable interconnect resources outside the block.

Four independent inputs are provided to each of two function generators (F1 - F4 and G1 - G4). These function generators, with outputs labeled F' and G', are each capable of implementing any arbitrarily defined Boolean function of four inputs. The function generators are implemented as memory look-up tables. The propagation delay is therefore independent of the function implemented.

A third function generator, labeled H', can implement any Boolean function of its three inputs. Two of these inputs can optionally be the F' and G' functional generator outputs. Alternatively, one or both of these inputs can come from outside the CLB (H2, H0). The third input must come from outside the block (H1). Signals from the function generators can exit the CLB on two outputs.

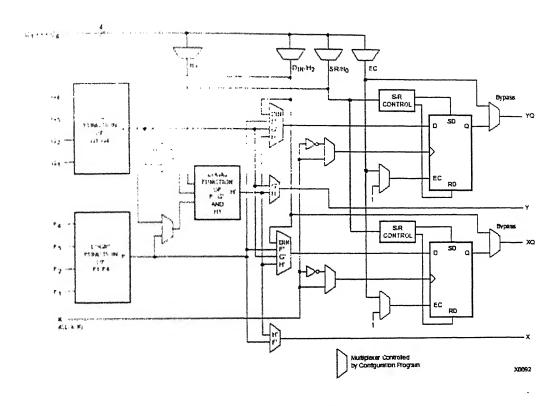


Fig. (1) Block Diagram of Xilinx 4000 series CLB

F' or H' can be connected to the X output. G' or H' can be connected to the Y output. A CLB can be used to implement any of the following functions:

- any function of up to four variables, plus any second function of up to four unrelated variables, plus any third function of up to three unrelated variables?
- any single function of five variables
- any function of four variables together with some functions of six variables
- some functions of up to nine variables.

Implementing wide functions in a single block reduces both the number of blocks required and the delay in the signal path, achieving both increased capacity and speed. The versatility of the CLB function generators significantly improves system speed. In addition, the design-software tools can deal with each function generator independently. This flexibility improves cell usage.

### Configuration Scheme:

The Xilinx FPGA uses the SRAM for the initialization, hence on power off the configuration will washout. So these FPGAs has to be reconfigure each time on power on. For the hardware using FPGA the configuration is done by PROM (crogrammable ROM). The choice of PROM depends on the target FPGA chosen for the design. The table (2) shows the requirement of the PROM for the implementation of tracking algorithms.

Device	Configuration Bits	PROM
XX 40 % F	147552	XC17128E
XC 4020 E	329312	XC1701

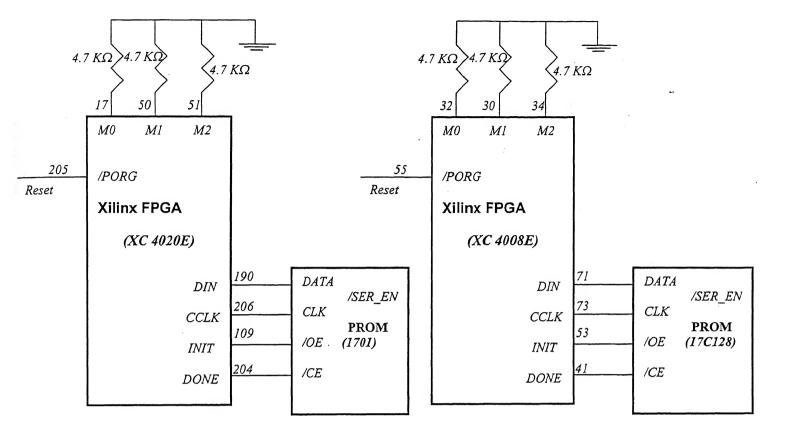
Table (2): PROM requirement for the implementation

There are various mode of operation for the FPGA reconfiguration. The table (3) gives the details for the mode selection and its requirement.

Mode	M2	M1	МО	CCLK	Data
Mester Seration	()	Ó	0	output	Bit-serial
Serve Series	1	1	1	input	Bit-serial
Master	ì	0	0	output	Bit-Wide
Parallel Up				·	Increment
, G. G op			_		From 0000
Master	1	1	0	output	Bit-Wide
Parallel Down					Decrement
, drailor bo					From FFFF
Peripheral	0	1	1	input	Bit-Wide
Synchronous					
Peripheral	1	0	1	output	Bit-Wide
Asynchronous					
Reserved	0	1	0		
Reserved	0	0	1		

Table (3): for Configuration Modes

The fig. (2) shows the configuration scheme for the proposed design. It uses the master-slave mode of configuration for FPGA.



Configuration Diagram for FPGA XC 4020 E

Configuration Diagram for FPGA XC 4008 E

Fig. 2: Configuration Scheme for the target FPGAs in Master-Slave Mode

#### Advantage of Re-configuration:

FPGA devices can be re-configured to change logic function while resident in the system. This capability gives the system designer a new degree of freedom not available with any other type of logic. Hardware can be changed as easily as software. Design updates or modifications are easy, and can be made to products already in the field. An FPGA can even be re-configured dynamically to perform different functions at different times. Re-configurable logic can be used to implement system self-diagnostics, create systems capable of being re-configured for different environments or operations, or implement multi-purpose hardware for a given application. As an added benefit, using reconfigurable FPGA devices simplifies hardware design and debugging and shortens product time-to-market.